

Processor Organization

Thoai Nam



Outline

- ❑ Criteria:
 - Diameter, bisection width, etc.
- ❑ Processor Organizations:
 - Mesh, binary tree, hypertree, pyramid, butterfly, hypercube, shuffle-exchange



Criteria

□ Diameter

- The largest distance between two nodes
- Lower diameter is better

□ Bisection width

The minimum number of edges that must be removed in order to divide the network into two halves (within one)

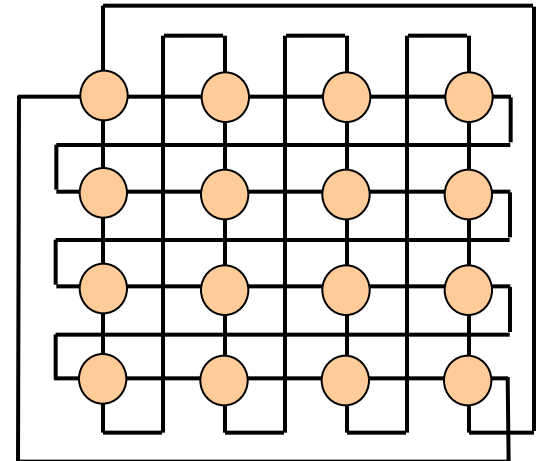
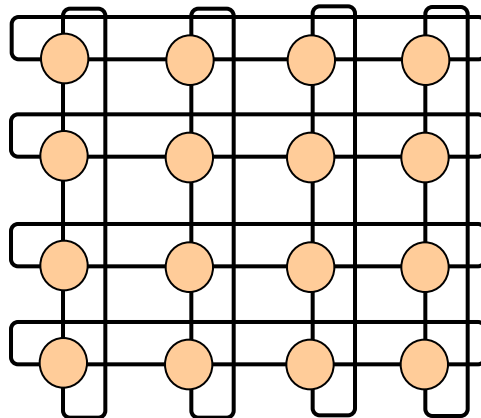
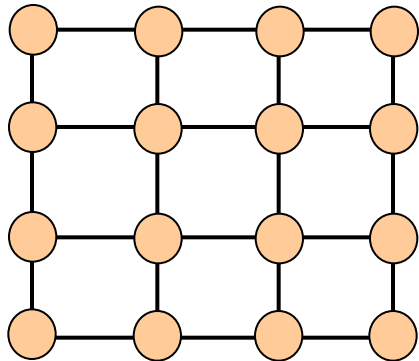
□ Number of edges per node

□ Maximum edge length



Mesh (1)

- Q-dimensional lattice
- Communication is allowed only between neighboring nodes. Interior nodes communicate with $2q$ other nodes.





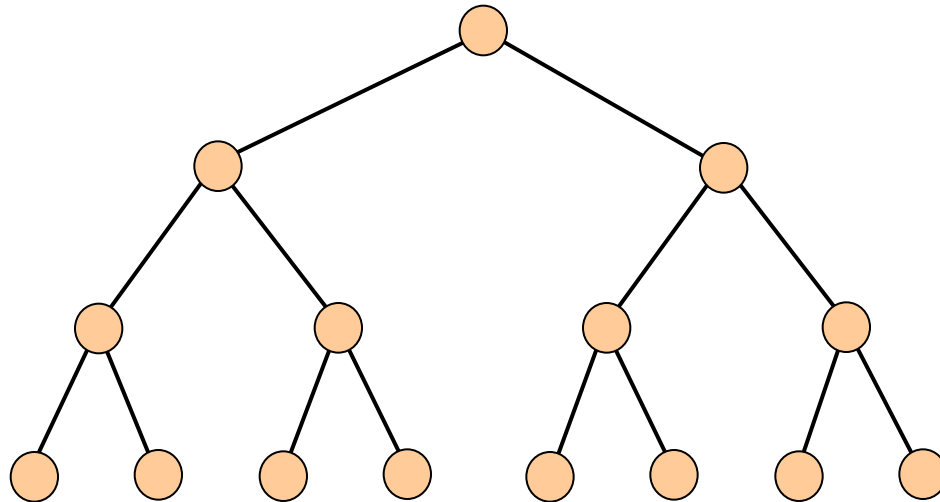
Mesh (2)

- Q-dimensional mesh with k^q nodes
 - Diameter: $q(k-1)$
 - Bisection width: k^{q-1}
 - The maximum number of edges per node: $2q$
 - The maximum edge length is a constant



Binary Tree

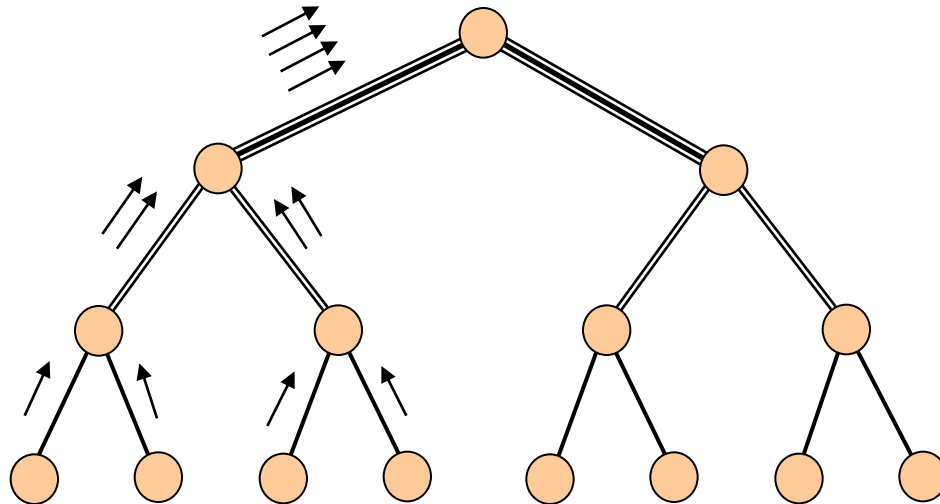
- Depth $k-1$: 2^k-1 nodes
- Diameter: $2(k-1)$
- Bisection width: 1
- Length of the longest edge: increasing





Fat Tree

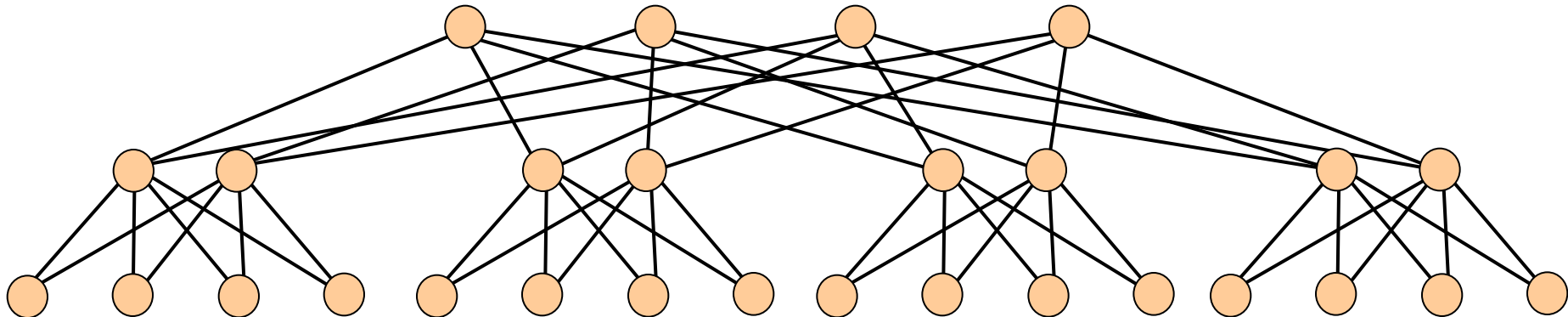
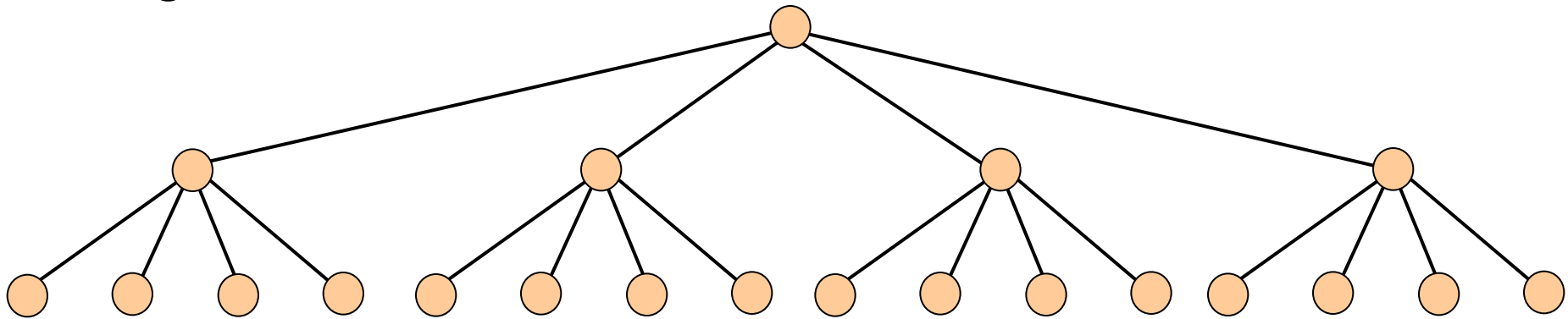
- Bandwidth problem on binary tree





Hypertree (1)

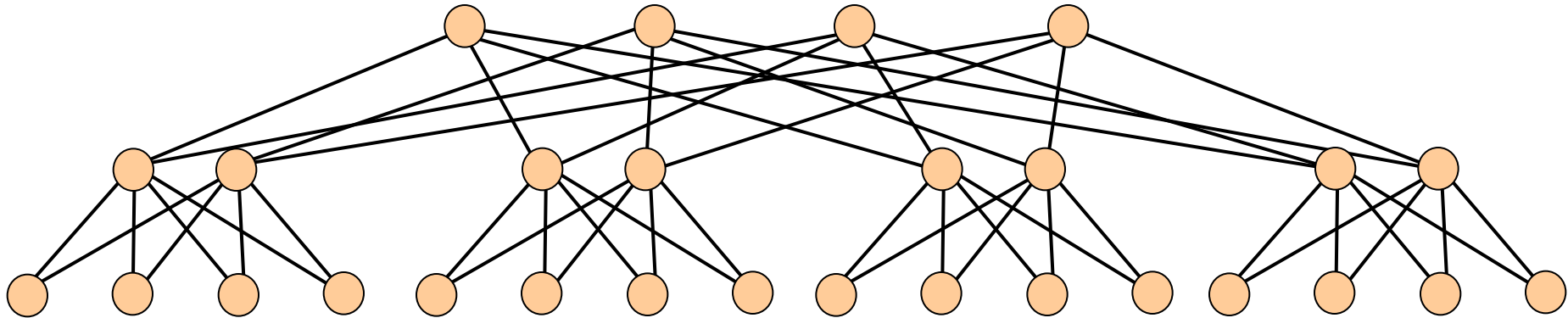
- Hypertree of **degree k** and **depth d**: a complete k-ary tree of height d.





Hypertree (2)

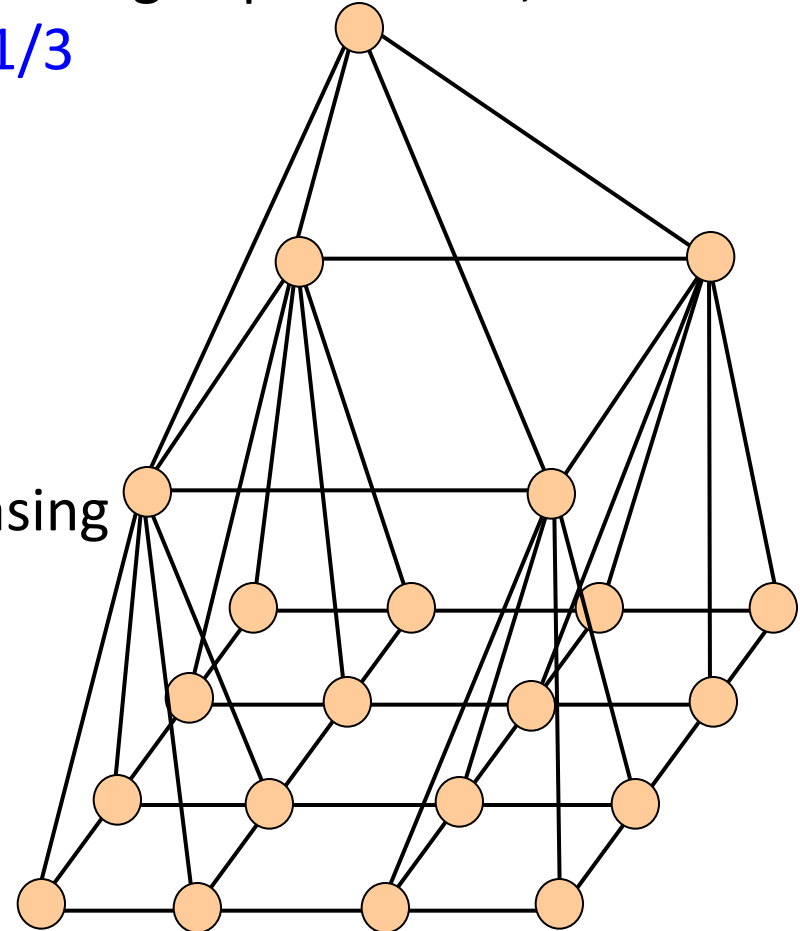
- A 4-ary hypertree with depth d has 4^d leaves and $2^d(2^{d+1}-1)$ nodes in all
 - Diameter: $2d$
 - Bisection width: 2^{d+1}
 - The number of edges per node ≤ 6
 - Length of the longest edge: increasing





Pyramid

- Size k^2 : base a 2D mesh network containing k^2 processors, the total number of processors= $(4/3)k^2 - 1/3$
- A pyramid of size k^2 :
 - Diameter: $2\log k$
 - Bisection width: $2k$
 - Maximum of links per node: 9
 - Length of the longest edge: increasing



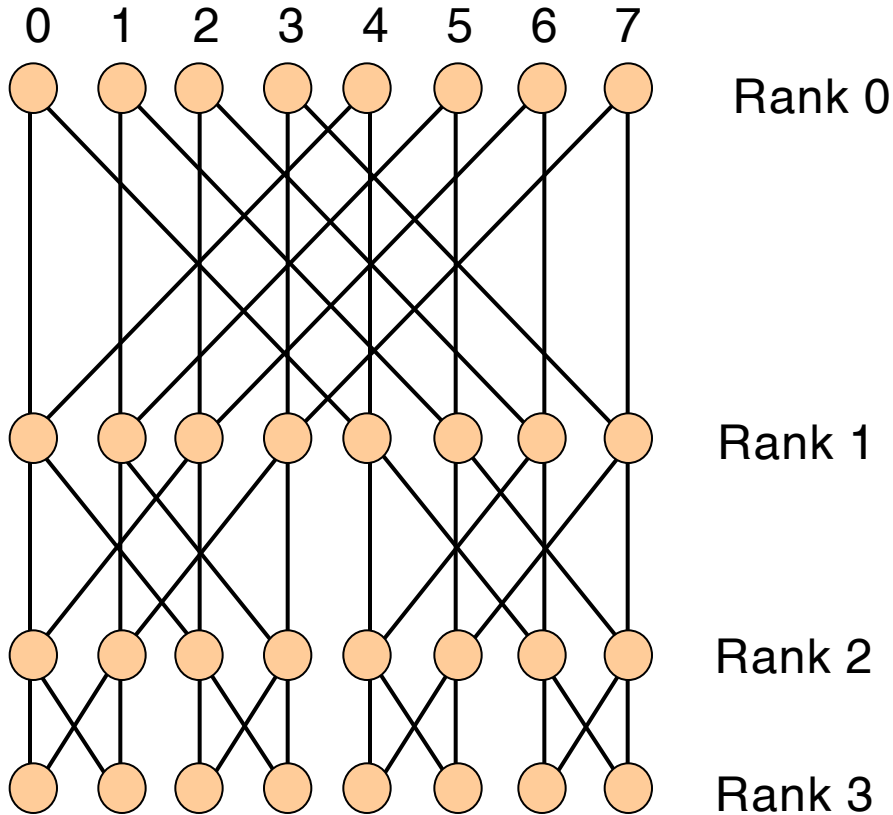


Butterfly (1)

- $(k+1)2^k$ nodes divided into $k+1$ rows (rank), each contains $n=2^k$ nodes.
- Ranks are labeled 0 through k
- **Node(i,j)**: j -th node on the i -th rank
- Node(i,j) is connected to two nodes on rank $i-1$: node($i-1,j$) and node ($i-1,m$), where m is the integer found by inverting the i -th most significant bit in the binary representation of j
- If **node(i,j)** is connected to **node($i-1,m$)**, then **node(i,m)** is connected to **node($i-1,j$)**
- Diameter= $2k$
- Bisection width= 2^k
- Length of the longest edge: increasing



Butterfly (2)



Node(1,5): $i=1, j=5$

$j = 5 = \mathbf{101}$ (binary)

$\downarrow i=1$

$\mathbf{001} = 1$

Node(1,5) is connected to
node(0,1)

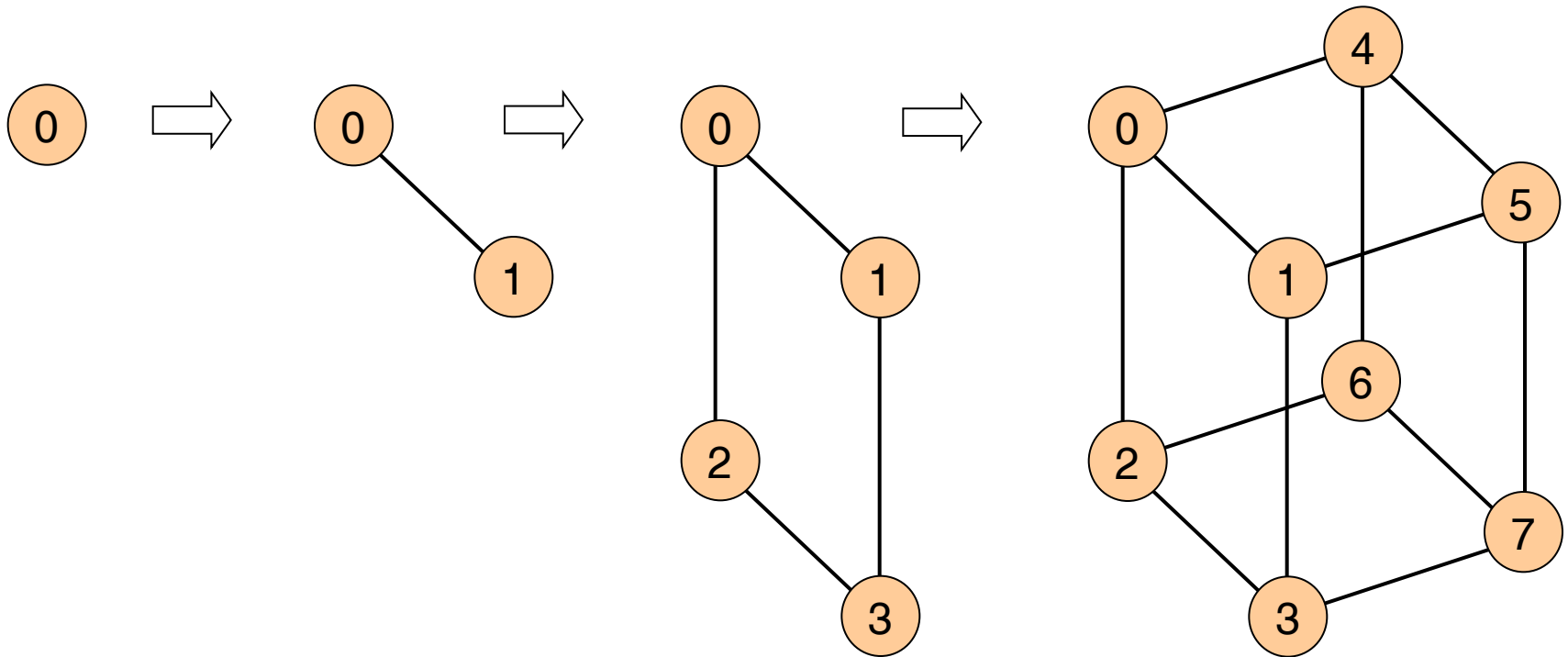


Hypercube (1)

- ❑ 2^k nodes form a k -dimensional hypercube
- ❑ Nodes are labeled $0, 1, 2, \dots, 2^k-1$
- ❑ Two nodes are adjacent if their labels differ in exactly one bit position
- ❑ Diameter= k
- ❑ Bisection width= 2^{k-1}
- ❑ Number of edges per node is k
- ❑ Length of the longest edge: increasing

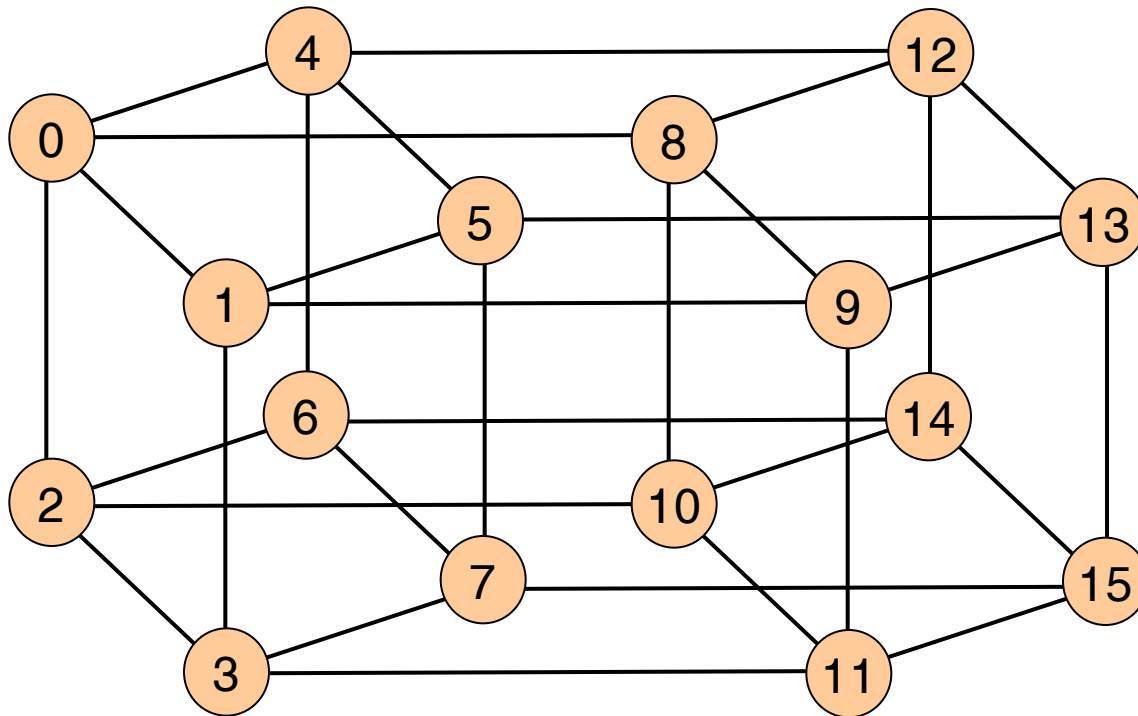


Hypercube (2)





Hypercube (3)



- 5 = **0101**
- 1 = **0001**
- 4 = **0100**
- 13 = **1101**



Others

- ❑ Torus

 - <http://clusterdesign.org/torus/>

 - <http://www.fujitsu.com/global/about/tech/k/whatis/network/>

- ❑ Cube-Connected cycles

- ❑ Shuffle-Exchange

- ❑ De Bruijn