

**Exercises**  
**(Course: Database Management Systems)**  
**Chapter 6**  
**Database Recovery Techniques**

1. Exercise 23.21 in the text book (“Fundamentals of Database Systems- 6th Edition”, Elmasri et al.)

Suppose that the system crashes before the [read\_item,  $T_3$ , A] entry is written to the log in Figure 23.1(b). Will that make any difference in the recovery process?

(a)

$T_1$
read_item(A)
read_item(D)
write_item(D)

$T_2$
read_item(B)
write_item(B)
read_item(D)
write_item(D)

$T_3$
read_item(C)
write_item(B)
read_item(A)
write_item(A)

**Figure 23.1**

Illustrating cascading rollback (a process that never occurs in strict or cascadeless schedules). (a) The read and write operations of three transactions. (b) System log at point of crash. (c) Operations before the crash.

(b)

	A	B	C	D
	30	15	40	20
	[start_transaction, $T_3$ ]			
	[read_item, $T_3$ , C]			
*	[write_item, $T_3$ , B, 15, 12]	12		
	[start_transaction, $T_2$ ]			
	[read_item, $T_2$ , B]			
**	[write_item, $T_2$ , B, 12, 18]	18		
	[start_transaction, $T_1$ ]			
	[read_item, $T_1$ , A]			
	[read_item, $T_1$ , D]			
	[write_item, $T_1$ , D, 20, 25]			25
	[read_item, $T_2$ , D]			
**	[write_item, $T_2$ , D, 25, 26]			26
	[read_item, $T_3$ , A]			

\*  $T_3$  is rolled back because it did not reach its commit point.

\*\*  $T_2$  is rolled back because it reads the value of item B written by  $T_3$ .

← System crash

2. Exercise 23.22 in the text book (“Fundamentals of Database Systems- 6th Edition”, Elmasri et al.)

Suppose that the system crashes before the [write\_item,  $T_2$ , D, 25, 26] entry is written to the log in Figure 23.1(b). Will that make any difference in the recovery process?

3. Exercise 23.23 in the text book (“Fundamentals of Database Systems- 6th Edition”, Elmasri et al.)

Figure 23.6 shows the log corresponding to a particular schedule at the point of a system crash for four transactions  $T_1$ ,  $T_2$ ,  $T_3$ , and  $T_4$ . Suppose that we use the *immediate update protocol* with checkpointing. Describe the recovery process from the system crash. Specify which transactions are rolled-back, which operations in the log are redone and which (if any) are undone, and whether any cascading rollback takes place.

[start_transaction, $T_1$ ]
[read_item, $T_1$ , $A$ ]
[read_item, $T_1$ , $D$ ]
[write_item, $T_1$ , $D$ , 20, 25]
[commit, $T_1$ ]
[checkpoint]
[start_transaction, $T_2$ ]
[read_item, $T_2$ , $B$ ]
[write_item, $T_2$ , $B$ , 12, 18]
[start_transaction, $T_4$ ]
[read_item, $T_4$ , $D$ ]
[write_item, $T_4$ , $D$ , 25, 15]
[start_transaction, $T_3$ ]
[write_item, $T_3$ , $C$ , 30, 40]
[read_item, $T_4$ , $A$ ]
[write_item, $T_4$ , $A$ , 30, 20]
[commit, $T_4$ ]
[read_item, $T_2$ , $D$ ]
[write_item, $T_2$ , $D$ , 15, 25]

← System crash

**Figure 23.6**

A sample schedule and its corresponding log.

4. Exercise 23.24 in the text book (“Fundamentals of Database Systems- 6th Edition”, Elmasri et al.)

Suppose that we use the deferred update protocol for the example in Figure 23.6. Show how the log would be different in the case of deferred update by removing the unnecessary log entries; then describe the recovery process, using your modified log. Assume that only REDO operations are applied, and specify which operations in the log are redone and which are ignored.

5. After a crash, we find the following log and transaction & dirty page tables at time of checkpoint:

**Table 1 – The log at point of crash**

LSN	LAST_LSN	TRAN_ID	TYPE	PAGE_ID	...
1	0	T1	update	C	...
2	1	T1	update	B	...

3	0	T2	update	C	...
4	begin_checkpoint				
5	end_checkpoint				
6	2	T1	commit		...
7	0	T3	update	A	...

*Table 1 – Transaction & Dirty Page Tables at time of checkpoint*

TRANSACTION TABLE			DIRTY PAGE TABLE	
TRANSACTION ID	LAST LSN	STATUS	PAGE ID	LSN
T1	2	in progress	C	1
T2	3	in progress	B	2

Specify the transaction and dirty page tables after the analysis phase.